

- Basic idea: $M \models M[G]$

- Forcing relation \Vdash - semantic (intuitive)
 \Vdash^* - syntactic (formal)

* Proof of equivalence is involved \Rightarrow you can first skip.

- Forcing theorem: $M[G] \models \phi \Rightarrow \exists p \in \mathbb{P}^M (p \Vdash \phi)$

* First skip $M[G] \models ZFC$

- Application of forcing: "adding κ -many new reals (functions $f \in 2^\omega$) to M "

Being def. in M means there is a $\Phi_\varphi(p, \bar{\tau}_1, \dots, \bar{\tau}_n)$

such that

$$\Phi_\varphi^M(p, \bar{\tau}_1, \dots, \bar{\tau}_n) \Leftrightarrow p \Vdash_{\mathbb{P}, M} \varphi(\bar{\tau}_1, \dots, \bar{\tau}_n)$$

\Rightarrow It is a (definable) class in M .

Use it:

$\mathbb{P} \in M$

Let $Y := \{p \in \mathbb{P} : p \Vdash \dots\}$

$X := \{ \tau \in \mathbb{Z} \subseteq \mathbb{P}^M : \exists p \Vdash \dots \tau \dots \}$

Then X, Y, \dots are $\in M$.

(because \Vdash is definable and $M \models$ Compr.)

$\mathbb{P} \in M$

$D \subseteq \mathbb{P}$ dense $D \in M$

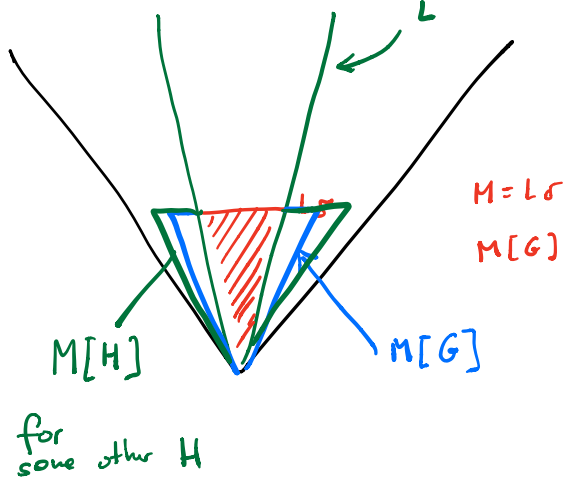
M transitive

$\Rightarrow D \in M$

$\mathbb{P} \in M$

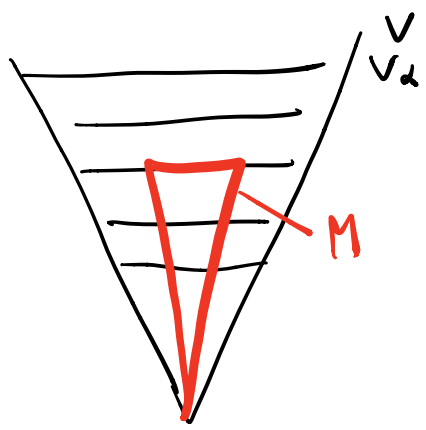
$G \subseteq \mathbb{P}$ will we $G \in M$

But $G \notin M$



$$o(M[G]) = o(M)$$

- $\text{Ord} \cap M$ is initial segm. of Ord
 $p \in \alpha \in M \rightarrow \beta \in M$
- $\text{Ord} \cap M$ is itself an ordinal $=: o(M)$



$x \in M$
 $\text{rk}(x) = \alpha \Rightarrow \alpha \in M$

	12:00	14:00	16:00
Wed 27	Talk 1	Talk 2	
Thu 28	Talk 3	Talk 4	Talk 5
Fr 29	Talk 6	Talk 7	

\curvearrowright Lidet + David