

## Proof that $M[G] \models$ Replacement.

Let  $a \in M[G]$  and  $\varphi$  formula.

leave out parameters  
for convenience.

Suppose  $M[G] \models \forall x \in a \exists y \varphi(x, y)$

To show: There is  $b \in M[G]$  such that

$$M[G] \models \forall x \in a \exists y \in b \varphi(x, y)$$

Let  $\tau$  be such that  $\tau_G = a$ .

Claim:  $\exists Q \in M[G]$  such that, for all  $\sigma \in \text{dom}(\tau)$  and  $p \in \mathbb{P}$

- If  $\exists \theta : p \Vdash \varphi(\sigma, \theta)$ , then  $\exists \theta \in Q : p \Vdash \varphi(\sigma, \theta)$

Proof of Claim: In  $M$ , define the following function:

$F(\sigma, p) :=$  1.) If  $\exists \theta$  s.t.  $p \Vdash \varphi(\sigma, \theta)$ , pick the one

with least rank to be the value

of  $F(\sigma, p)$

2.) Otherwise,  $\emptyset$ .

This argument  
is often used, sometimes  
called "Scott's Trick".

Then  $F$  is welldefined in  $M$ . Moreover  $\text{dom}(\tau) \times \mathbb{P} \in M$ ,

and  $F[\text{dom}(\tau) \times \mathbb{P}] \in Q$ . Since  $M \models \text{REPL}$ ,  $Q \in M$

$\square$  (Claim)

Now, let  $\pi := Q \times \{1\}$ . Then it is not hard to verify:

$$M[G] \models \forall x \in \tau_G \exists y \in \pi_G \varphi(x, y). \quad \square$$