

## ANALYTIC DETERMINACY

**Notation.** Fix a bijection  $i \mapsto s_i$  from  $\omega \rightarrow \omega^{<\omega}$  such that if  $s_i \subseteq s_j$ , then  $i \leq j$ . (This implies that  $\text{lh}(s_i) \leq i$ .) Let  $T \subseteq (\omega \times \omega)^{<\omega}$  be a tree,  $x \in \omega^\omega$ , and  $s \in \omega^{<\omega}$ . Then we let

$$\begin{aligned} T_s &:= \{t \in \omega^{<\omega}; (s \upharpoonright \text{lh}(t), t) \in T\}, \\ T_x &:= \{t \in \omega^{<\omega}; (x \upharpoonright \text{lh}(t), t) \in T\} = \bigcup_{n \in \mathbb{N}} T_{x \upharpoonright n}, \\ K_s &:= \{i \leq \text{lh}(s); s_i \in T_s\}, \text{ and} \\ K_x &:= \{i \in \omega; s_i \in T_x\} = \bigcup_{n \in \mathbb{N}} K_{x \upharpoonright n}. \end{aligned}$$

We note that  $T_s$  is a tree of finite height (every element  $t \in T_s$  has length  $\leq \text{lh}(s)$ ) and that  $K_s$  is a finite set. We observe that  $T_x = \{s_i; i \in K_x\}$  (but, in general,  $T_s \not\supseteq \{s_i; i \in K_s\}$ ).

We remember that if  $A \in \mathbf{\Pi}_1^1$ , then there is a tree  $T$  on  $\omega \times \omega$  such that

$$\begin{aligned} x \in A &\text{ if and only if } (T_x, \supseteq) \text{ is wellfounded} \\ &\text{ if and only if } (T_x, <_{\text{KB}}) \text{ is wellordered} \\ &\text{ if and only if there is an order preserving map from } (T_x, <_{\text{KB}}) \text{ to } (\omega_1, <) \end{aligned}$$

where  $<_{\text{KB}}$  is the Kleene-Brouwer order on  $\omega^{<\omega}$ . For any  $s \in \omega^{<\omega}$ , we write  $<_s$  for the order induced by the Kleene-Brouwer order on  $\omega^{<\omega}$  on  $K_s$ , i.e.,  $i <_s j$  if and only if  $s_i <_{\text{KB}} s_j$ . Note that since  $K_s$  is finite,  $(K_s, <_s)$  is a (finite) wellorder.

Let  $S$  be any tree on  $\omega$  and  $\kappa$  be an uncountable cardinal. A function  $g : \omega \rightarrow \kappa$  is called a *KB-code* for  $S$  if for all  $i$  and  $j$  such that  $s_i, s_j \in S$ , we have that  $s_i <_{\text{KB}} s_j \leftrightarrow g(i) < g(j)$ . Clearly, there is an order preserving map from  $(S, <_{\text{KB}})$  to  $(\omega_1, <)$  if and only if there is a KB-code for  $S$ , so we can add the following equivalence to the above characterisation of  $\mathbf{\Pi}_1^1$  sets:

$$x \in A \text{ if and only if there is a KB-code for } T_x$$

**Shoenfield's Theorem.** We first prove a tree representation theorem for  $\mathbf{\Pi}_1^1$  sets.

**Theorem 1** (Shoenfield). If  $\kappa$  is uncountable, then every  $\mathbf{\Pi}_1^1$  set is  $\kappa$ -Suslin.

*Proof.* Let  $A \in \mathbf{\Pi}_1^1$  and let  $T$  be a tree on  $\omega \times \omega$  such that  $x \in A$  if and only if there is a KB-code for  $T_x$ . Let  $M$  be the set of all partial functions from  $\omega$  into  $\kappa$  with finite domain. Note that  $|M| = \kappa$ , so it is sufficient to show that  $A$  is  $M$ -Suslin. If  $s \in \omega^{<\omega}$  and  $u \in M^{<\omega}$  such that  $\text{lh}(h) \leq \text{lh}(s)$ , we say that  $u$  is *coherent with  $s$*  if

- (1) for all  $i < \text{lh}(u)$ , we have that  $\text{dom}(u_i) = K_{s \upharpoonright i}$ ,
- (2) for all  $i < \text{lh}(u)$ ,  $u(i)$  is an order preserving map from  $(K_{s \upharpoonright i}, <_{\text{KB}})$  into  $(\kappa, <)$ , and
- (3) for  $i \leq j$ , we have that  $u_i \subseteq u_j$ .

We now define the *Shoenfield tree* on  $\omega \times M$  by  $\widehat{T} := \{(s, u); u \text{ is coherent with } s\}$  and claim that  $A = \text{p}[\widehat{T}]$ :

“ $\subseteq$ ”: If  $x \in A$ , then let  $g : \omega \rightarrow \kappa$  be a KB-code for  $T_x$  and define  $u(i) := g \upharpoonright K_{x \upharpoonright i}$ . By definition,  $u \upharpoonright n$  is coherent with  $x \upharpoonright n$  for all  $n$ , and so  $(x, u) \in [\widehat{T}]$ .

“ $\supseteq$ ”: If  $x \in \text{p}[\widehat{T}]$ , find  $u \in M^\omega$  such that  $(x, u) \in [\widehat{T}]$ ; this means that for each  $n$ ,  $u \upharpoonright n$  is coherent with  $x \upharpoonright n$ . As noted above, we have that  $T_x = \{s_i; i \in K_x\} = \{s_i; \exists n (i \in \text{dom}(u(n)))\}$ . We define  $\widehat{u} := \bigcup \{u(i); i \in \omega\}$ . By coherence,  $\widehat{u}$  is a function from  $K_x$  to  $\kappa$ ; now we define

$$g : \omega \rightarrow \kappa : n \mapsto \begin{cases} \widehat{u}(n) & \text{if } n \in K_x \text{ and} \\ 0 & \text{otherwise.} \end{cases}$$

We claim that  $g$  is a KB-code for  $T_x$  whence  $x \in A$ : Suppose not, then there are  $i$  and  $j$  such that  $s_i, s_j \in T_x$  and  $s_i <_{\text{KB}} s_j \not\leftrightarrow g(i) < g(j)$ . Since  $i, j \in K_x$ , find  $n$  large enough such that  $i, j \in K_{x \upharpoonright n}$ . By definition  $g \upharpoonright K_{x \upharpoonright n} = u(n)$ . But this means that  $u(n)$  is not an order preserving map from  $(K_{x \upharpoonright n}, <_{\text{KB}})$  into  $(\kappa, <)$ , violating condition (3) of coherence. Q.E.D

**Measurable Cardinals.** Let  $X$  be a set. A non-empty family  $U \subseteq \wp(X)$  is called a *ultrafilter over  $X$*  if for any  $A, B \subseteq X$ , we have that

- (1) if  $A, B \in U$ , then  $A \cap B \in U$ ,
- (2) if  $A \in U$  and  $B \supseteq A$ , then  $B \in U$ , and
- (3) if  $A \notin U$ , then  $X \setminus A \in U$ .

We say that an ultrafilter is *non-trivial* if it does not contain any finite sets and if  $\kappa$  is any cardinal, it is called  *$\kappa$ -complete* if it is closed under intersections of size  $< \kappa$ . Note that  $\omega$ -completeness follows from (1). A non-trivial  $\kappa$ -complete ultrafilter cannot contain any sets of size  $< \kappa$ .

[If  $|A| = \lambda < \kappa$ , then for each  $a \in A$ ,  $\{a\} \notin U$ , so by (3),  $X \setminus \{a\} \in U$ , but then by  $\kappa$ -completeness,  $X \setminus A = \bigcap \{X \setminus \{a\} ; a \in A\} \in U$ . If now  $A \in U$ , then  $\emptyset = A \cap X \setminus A \in U$ . Contradiction to non-triviality.]

An uncountable cardinal  $\kappa$  is called *measurable* if there is a  $\kappa$ -complete non-trivial ultrafilter on  $\kappa$ . The Axiom of Choice implies that there are non-trivial ultrafilters on  $\omega$ ; as mentioned, they are  $\omega$ -complete, so  $\aleph_0$  technically satisfies the conditions of the definition. The existence of uncountable measurable cardinals cannot be proved in ZFC and is a so-called *large cardinal axiom*. More precisely, if MC stands for “there is a measurable cardinal”, then for every model  $M \models \text{ZFC} + \text{MC}$ , I can find a submodel  $N \subseteq M$  such that  $N \models \text{ZFC} + \neg \text{MC}$ .

Being measurable has interesting consequences for the combinatorics on  $\kappa$ . We are going to use one of them in our proof of analytic determinacy. As usual, we denote by  $[\kappa]^n$  the set of  $n$ -element subsets of  $\kappa$ . A function  $f : [\kappa]^n \rightarrow \omega$  is called an  *$n$ -colouring* and a set  $H$  is called *homogeneous for  $f$*  if  $f \upharpoonright [H]^n$  is constant. We call  $f$  a *finite colouring* if it is an  $n$ -colouring for some natural number  $n \in \mathbb{N}$ .

**Theorem 2** (Rowbottom). If  $\kappa$  is measurable, then for every countable set  $\{f_s ; s \in S\}$  of finite colourings, there is a set  $H$  of size  $\kappa$  that is homogeneous for all colourings  $f_s$ .

In our proof of analytic determinacy, we are only going to use Rowbottom’s Theorem, no other properties of measurable cardinals; so, for our purposes, one could take the statement of Rowbottom’s Theorem as the assumption for analytic determinacy in the next section.

**Analytic Determinacy.** If  $\Gamma$  is a boldface pointclass, then  $\text{Det}(\Gamma)$  is equivalent to  $\text{Det}(\check{\Gamma})$ . Thus, analytic determinacy and co-analytic determinacy are equivalent.

**Theorem 3** (Martin, 1969/70). If there is a measurable cardinal, then every co-analytic set is determined.

**Proof.** Let  $\kappa$  be a measurable cardinal and  $A \in \Pi_1^1$ . We aim to show that the game  $G(A)$  is determined. By (the proof of) Shoenfield’s Theorem, we know that there is a tree  $\widehat{T}$  on  $\omega \times M$  such that  $A = p[\widehat{T}]$ . (Remember that  $M$  was the set of partial functions from  $\omega$  to  $\kappa$  with finite domain.) We are going to define a (determined) game  $G_{\text{aux}}(\widehat{T})$  based on the Shoenfield tree and show that a winning strategy for either player in  $G_{\text{aux}}(\widehat{T})$  can be transformed into a winning strategy for the same player in the *original game*  $G(A)$ . This proves the theorem.

In the *auxiliary game*, player I plays elements of  $\omega \times M$  and player II plays elements of  $\omega$  as follows:

$$\begin{array}{c|cccccc} \text{I} & x_0, u_0 & x_2, u_1 & x_4, u_2 & x_6, u_3 & \cdots \\ \text{II} & & x_1 & x_3 & x_5 & x_7 \cdots \end{array}$$

We obtain a sequence  $x \in \omega^\omega$  with  $x(n) := x_n$  and a sequence  $u \in M^\omega$  with  $u(n) := u_n$ . Player I wins  $G_{\text{aux}}(\widehat{T})$  if  $(x, u) \in [\widehat{T}]$ . Note that  $G_{\text{aux}}(\widehat{T})$  is a closed game on  $\omega \times M$ , thus by the Gale-Stewart Theorem, it is determined.

Let us make a number of observations about the relationship between the original game  $G(A)$  and the auxiliary game  $G_{\text{aux}}(\widehat{T})$ . We call the moves  $u_i$  *auxiliary moves*. If  $p$  is a position in the auxiliary game (i.e., a finite sequence of elements of  $\omega$  and elements of  $M$  in the right order), then we can define a position  $p^*$

in the original game by forgetting about the auxiliary moves. This allows us to consider strategies  $\tau$  for player II in the original game as strategies in the auxiliary game: if  $p$  is a position in the auxiliary game, we let  $\tau_*(p) := \tau(p^*)$ , i.e., just forget about the auxiliary moves and play as if you were playing in the original game.

**Lemma 4.** If player I has a winning strategy in  $G_{\text{aux}}(\widehat{T})$ , then they have a winning strategy in  $G(A)$ .

*Proof.* Suppose  $\sigma$  is a winning strategy in  $G_{\text{aux}}(\widehat{T})$  and  $\tau$  is any strategy for player II in the original game. As just mentioned, then  $\tau_*$  is the version of that strategy in  $G_{\text{aux}}(\widehat{T})$ . Since  $\sigma$  is winning, we know that  $\sigma * \tau_* = (x, u) \in [\widehat{T}]$ . Define a strategy  $\sigma^*$  in the original game as follows: while player II plays natural number moves according to  $\tau$ , you produce the auxiliary play  $\sigma * \tau_*$  on an auxiliary board. If that auxiliary game tells you to produce a position  $p$  by your next move, then you produce the move  $p^*$  in the original game. Then  $\sigma^* * \tau = x$ , and thus  $x \in p[\widehat{T}] = A$ , so  $\sigma^*$  is winning. Q.E.D

**Lemma 5.** If player II has a winning strategy in  $G_{\text{aux}}(\widehat{T})$ , then they have a winning strategy in  $G(A)$ .

*Proof.* Let  $s \in \omega^{<\omega}$ . Let  $k_s := |K_s|$ . If  $Q \in [\kappa]^{k_s}$ , then there is a unique order preserving map  $w : (K_s, <_s) \rightarrow (Q, <)$ . Let  $u_i^{s,Q} := w \upharpoonright K_{s \upharpoonright i}$ . Then  $(u_i^{s,Q}; i < \text{lh}(s))$  is coherent with  $s$ . Thus, if you fix some  $Q \in [\kappa]^{k_s}$ , you can transform a position  $s$  in the original game into a position  $s_{*,Q}$  in the auxiliary game in such a way that the auxiliary moves produce  $Q$  as the range and form a sequence coherent with the position  $s$ .

Let now  $\tau$  be a strategy for player II in the auxiliary game. For each  $s \in \omega^{<\omega}$ , we define a  $k_s$ -colouring  $f_s : [\kappa]^{k_s} \rightarrow \omega$  by  $f_s(Q) := \tau(s_{*,Q})$ : we colour the  $k_s$ -element subsets of  $\kappa$  by the answer that the strategy  $\tau$  gives to the position  $s$  augmented via  $Q$  in the sense given above. By Rowbottom's theorem, there is a set  $H \subseteq \kappa$  of size  $\kappa$  that is homogeneous for all functions  $f_s$ , i.e., if  $Q, Q' \in [H]^{k_s}$ , then  $\tau(s_{*,Q}) = f_s(Q) = f_s(Q') = \tau(s_{*,Q'})$ , so the answer of the strategy  $\tau$  does not depend on the set  $Q$  as long as it is a subset of  $H$ . In particular, we can take the simplest imaginable subset of  $H$  with  $k_s$  elements: let  $Q_{H,s}$  be the set consisting of the first  $k_s$  many elements of  $H$ .

Now, we define a strategy  $\tau_H$  for player II in the original game by  $\tau_H(s) := \tau(s_{*,Q_{H,s}})$ . (Note that the precise choice of the set  $Q_{H,s}$  is irrelevant in this definition by homogeneity, since  $f_s(Q_{H,s}) = f_s(Q)$  for any  $Q \in [H]^{k_s}$ .)

We prove that if  $\tau$  was winning in the auxiliary game, then  $\tau_H$  is winning in the original game. Suppose not, so there is a counterstrategy  $\sigma$  such that  $x := \sigma * \tau_H \in A$ . This means (since  $H$  is uncountable) that there is an orderpreserving map from  $(T_x, <_{\text{KB}})$  to  $(H, <)$  giving rise to a KB-code  $g : \omega \rightarrow H$  for  $T_x$ . Using the KB-code  $g$ , we can now define  $u_i := g \upharpoonright K_{x \upharpoonright i}$  and consider the play of the auxiliary game

$$\begin{array}{c|cccccc} \text{I} & x_0, u_0 & x_2, u_1 & x_4, u_2 & x_6, u_3 & \cdots \\ \text{II} & x_1 & x_3 & x_5 & x_7 & \cdots \end{array}$$

producing  $(x, u) \in [\widehat{T}]$ . We claim that this is a play according to  $\tau$ , so we need to show that for every  $i \in \mathbb{N}$ , the play by player II is the  $\tau$ -answer to the previous position, i.e.,  $x_{2i+1} = \tau(x_0, u_0, x_1, \dots, x_{2i}, u_i)$ . Fix  $i \in \mathbb{N}$  and consider  $Q := \text{ran}(u_i) \subseteq H$ . Then we have that  $(x_0, u_0, x_1, \dots, x_{2i}, u_i) = (x \upharpoonright 2i + 1)_{*,Q}$ . We see that

$$\begin{aligned} x_{2i+1} &= \tau_H(x \upharpoonright 2i + 1) && \text{(since } x \text{ was produced by } \tau_H) \\ &= \tau((x \upharpoonright 2n + 1)_{*,Q_{H,s}}) && \text{(by definition of } \tau_H) \\ &= \tau((x \upharpoonright 2n + 1)_{*,Q}) && \text{(since the choice of } Q \text{ doesn't matter by homogeneity)} \\ &= \tau(x_0, u_0, x_1, \dots, x_{2i}, u_i), \end{aligned}$$

so the above play is a play according to  $\tau$ . But that is a contradiction, since  $\tau$  was winning for player II, and so  $(x, u) \notin [\widehat{T}]$ . Q.E.D

Lemmas 4 & 5 together with the fact that  $G_{\text{aux}}(\widehat{T})$  was determined (since it is a closed game) imply that  $G(A)$  is determined. Q.E.D