Mackey and Tambara functors

Birgit Richter

13th of January 2025

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$$\begin{array}{ccc}
U & \stackrel{\alpha}{\longrightarrow} & V \\
\downarrow^{\beta} & & \downarrow^{\gamma} \\
W & \stackrel{\delta}{\longrightarrow} & Z
\end{array}$$

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▶ for every pair of finite *G*-sets *X* and *Y*, applying M_* to $X \to X \sqcup Y \leftarrow Y$ gives the component maps of an isomorphism $\underline{M}(X) \oplus \underline{M}(Y) \cong \underline{M}(X \sqcup Y)$.

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The covariant part encodes transfer maps: For K < H < G and the canonical projection $p: G/K \to G/H$ we get a *transfer map* $M_*(p) = \operatorname{tr}_K^H: \underline{M}(G/K) \to \underline{M}(G/H).$

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Example Let *B* be an abelian group with a *G*-action. Then the fixed point Mackey functor \underline{B}^{fix} has $\underline{B}^{fix}(G/H) = B^H \cong G$ -maps(G/H, B).

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$$\operatorname{res}_{K}^{H} := (B^{fix})^{*}(p) \colon B^{H} \to B^{K}.$$

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The transfer $\operatorname{tr}_{K}^{H}$ for $p: G/K \to G/H$ sends an $f \in G$ -maps(G/K, B) to $\operatorname{tr}_{K}^{H}(f)(gH) = \sum_{x \in p^{-1}(gH)} f(x)$.

Definition The Lindner category, B_G^+ , has as objects finite G-sets.

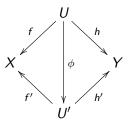
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Here the above span is equivalent to $X \stackrel{f'}{\longleftarrow} U' \stackrel{h'}{\longrightarrow} Y$, if there is a bijection of finite *G*-sets $\phi: U \rightarrow U'$ such that

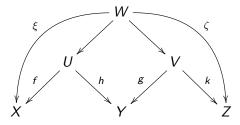


commutes.

Composition in B_G^+ is defined via pullbacks:

$$[Y \stackrel{g}{\longleftrightarrow} V \stackrel{k}{\longrightarrow} Z] \circ [X \stackrel{f}{\longleftrightarrow} U \stackrel{h}{\longrightarrow} Y] := [X \stackrel{\xi}{\longleftrightarrow} W \stackrel{\zeta}{\longrightarrow} Z]$$

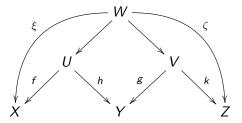
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The set $B_G^+(X, Y)$ carries an abelian monoid structure defined via

$$[X \xleftarrow{f_1} U_1 \xrightarrow{h_1} Y] + [X \xleftarrow{f_2} U_2 \xrightarrow{h_2} Y] := [X \xleftarrow{(f_1, f_2)} U_1 \sqcup U_2 \xrightarrow{(h_1, h_2)} Y]$$

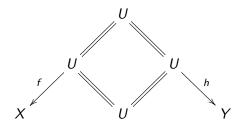
A Mackey functor can then be described as an additive functor $\underline{M}: B_{\mathcal{G}} \rightarrow Ab$.

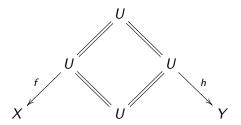
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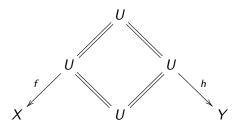
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The equivalence of the two definitions can be seen as follows: For a map $f: X \to Y$ of finite *G*-sets we consider the span $X = X \xrightarrow{f} Y$ which gives $M_*(f)$, and the span $Y \xleftarrow{f} X = X$ which yields $M^*(f)$.



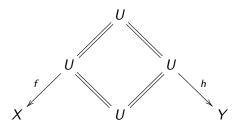


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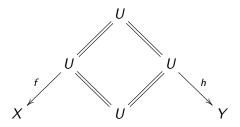
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► The Burnside Mackey functor, <u>A</u>, for G sends a finite G-set X to the Grothendieck group of the abelian monoid of isomorphism classes of finite G-sets over X.

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Later, we'll need to change the group: If H < G, then we can restrict a *G*-Mackey functor \underline{M} to an *H*-Mackey functor. We'll denote this by $i_{H}^{*}(\underline{M})$. Explicitly: $i_{H}^{*}(\underline{M})(T) := \underline{M}(G \times_{H} T)$ for any *H*-set *T*. Here, two classes $[Y \to X]$ and $[Z \to X]$ are added to give $[Y \sqcup Z \to X]$. A map $f: X \to W$ of finite *G*-sets induces $A_*(f): \underline{A}(X) \to \underline{A}(W)$

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We can also induce an *H*-Mackey functor \underline{N} up to a *G*-Mackey functor by sending a *G*-set *S* to $\operatorname{ind}_{H}^{G}\underline{N}(S) := \underline{N}(S)$ where we view *S* just as an *H*-set.

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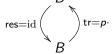
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For the C_2 -fixed point Mackey functor of $\mathbb{Z}[i]$ the Lewis diagram is

$$\mathbb{Z} = \mathbb{Z}[i]^{C_2}$$

res=inclusion $\left({
ight)} \operatorname{tr}(1)=2, \operatorname{tr}(i)=i-i=0$ $\mathbb{Z}[i]$

 $W_{C_2}(e) = C_2, 1 \mapsto 1, i \mapsto -i.$

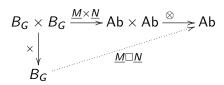
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$$\begin{array}{c|c} B_G \times B_G \xrightarrow{\underline{M} \times \underline{N}} Ab \times Ab \xrightarrow{\otimes} Ab \\ \times & & \\ B_G & \underline{M} \square \underline{N} \end{array}$$

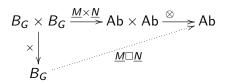
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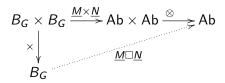


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This is also called a *Day convolution product* after Day, who defined this is his thesis in 1970: You merge the symmetric monoidal structures \times for finite *G*-sets and \otimes for abelian groups to get \Box for the functor category.

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Let's unravel that: The C_p -quotient is wrt the coordinatewise action. We denote the equivalence class of $x \otimes y$ wrt that action by $[x \otimes y]$. The transfer sends $m \otimes n \in \underline{M}(C_p/e) \otimes \underline{N}(C_p/e)$ to $[m \otimes n] \in (\underline{M}(C_p/e) \otimes \underline{N}(C_p/e))/C_p$.

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Example Let's take the *G*-Burnside Mackey functor with $\underline{A}(X) \cong B_G(G/G, X)$. As $G/G \times Y \cong Y$ for all finite *G*-sets *Y* and by abstract non-sense about representable functors and Day convolution products, we get that \underline{A} is the unit for \Box :

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Frobinius-Reciprocity (FR) identifies:

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This is a lot of structure...

Example Let's take the *G*-Burnside Mackey functor with $\underline{A}(X) \cong B_G(G/G, X)$. As $G/G \times Y \cong Y$ for all finite *G*-sets *Y* and by abstract non-sense about representable functors and Day convolution products, we get that \underline{A} is the unit for \Box : For all $\underline{M} \in G$ -Mack

$$\underline{A} \Box \underline{M} \cong \underline{M} \cong \underline{M} \Box \underline{A}.$$

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Commutative monoids in $(G-Mack, \Box)$ are commutative Green functors. They *cannot* be tensored with finite *G*-sets in a functorial manner. In order to have for instance maps $G/K \otimes \underline{R} \rightarrow G/H \otimes \underline{R}$ for K < H < G we need more structure. Tambara functors are Mackey functors with an additional multiplicative structure and with multiplicative norms:

For the map $p: G/K \to G/H$ we have a multiplicative map $N_p: \underline{R}(G/K) \to \underline{R}(G/H).$

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<u>*A*</u> is initial in Tamb_{*G*} with the product induced by the product of finite *G*-sets.

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For a map of finite *G*-sets $f: X \to Y$ we consider the pullback functor that sends $h: B \to Y$ to $X \times_Y B \to X$. This is a functor from the category of finite *G*-sets over *Y* to finite *G*-sets over *X*.

$$\prod_{f} A := \{(y,s) \mid y \in Y, s \colon f^{-1}(y) \to A \text{ with } ps(x) = x \forall x \in f^{-1}(y)\}.$$

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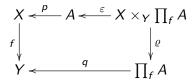
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A diagram isomorphic to



is called an *exponential diagram*. (Here, $\varepsilon(x, (y, s)) = s(x)$, $\varrho(x, (y, s)) = (y, s)$.)

Definition The category \mathcal{P}^{G} has as objects finite *G*-sets

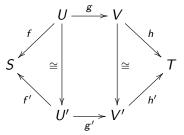
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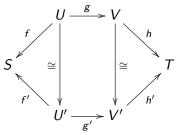
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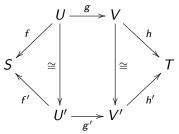


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Composition of morphisms is a bit involved. We define restriction, norm and transfer associated to a map:

Let $f: S \to T$ be a map of finite *G*-sets.

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 R is a contravariant functor from finite G-sets to P^G and N, T are covariant ones. Let $f: S \to T$ be a map of finite *G*-sets.

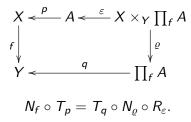
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- For a pullback diagram in finite G sets



$$R_g \circ N_f = N_{f'} \circ R_{g'}$$
 and $R_g \circ T_f = T_{f'} \circ R_{g'}$





Theorem [Kristen Mazur 2013, Rolf Hoyer 2014] There is a functor

$$(-)\otimes (-)\colon G\operatorname{-Sets}^{\mathsf{f}} \times \operatorname{Tamb}_{G} \to \operatorname{Tamb}_{G}$$

 $(X, R) \mapsto X \otimes R$

which satisfies the following properties:

Theorem [Kristen Mazur 2013, Rolf Hoyer 2014] There is a functor

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which satisfies the following properties:

1. For all X and Y in G-Sets^f and <u>R</u>, <u>T</u> in Tamb_G, there are natural isomorphisms $(X \sqcup Y) \otimes \underline{R} \cong (X \otimes \underline{R}) \Box (Y \otimes \underline{R})$

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- 2. There is a natural isomorphism $X \otimes (Y \otimes \underline{R}) \cong (X \times Y) \otimes \underline{R}$.
- 3. On the category with objects finite sets with trivial *G*-action and morphisms consisting only of isomorphisms, the functor restricts to exponentiation $X \otimes \underline{R} = \prod_{x \in X} \underline{R}$.

Definition Let G be a finite group, $\underline{R} \in \text{Tamb}_G$ and let X be a finite simplicial G-set.

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